

Addendum to the CM-5 C* User's Guide

Version 7.1, May 1993

1 Physical I/O Routines

Section 2.4.3 of the *CM-5 C* User's Guide* omits a discussion of routines for performing physical I/O; these routines are available in Version 7.1.

The physical I/O routines are `CMFS_write_file_physical` and `CMFS_read_file_physical`. They have the same interfaces as `CMFS_write_file_always` and `CMFS_read_file_always`. Namely:

```
int CMFS_write_file_physical(int fd,
                             void* source,
                             int bytes_per_position)

int CMFS_read_file_physical(int fd,
                            void* destination,
                            int bytes_per_position)
```

These routines are faster than the corresponding standard I/O routines and use less memory. However, note these points about their use:

- They are not portable. If you use the `CMFS_write_file_physical` routine to write data to an external storage device, you can recover the data only by using the `CMFS_read_file_physical` routine to read the data onto the same set of physical processors, using the same file system, and into a parallel variable with the same memory layout as the parallel variable from which the original write occurred. Essentially, you must use the same partition size, file system, and shape for both writing and reading; if you do a physical write, you must do a physical read.

- These routines return the total number of bytes read or written, rather than the number of bytes read or written per-position.
- The I/O data read or written by these routines is padded. When the file pointer is not located at an alignment boundary appropriate to the I/O device where the file referenced through the `fd` argument resides, the file pointer is adjusted upward to such a boundary before the transfer of the data between the CM and the I/O device begins. Furthermore, garbage data is written to pad the length of the I/O transaction to the same alignment. The alignment padding used is 512 bytes for Datavault IO, and 16 bytes for SDA IO.

To anticipate the impact of padding before doing a read or write, use this routine:

```
int CMFS_physical_transfer_length(int fd,
                                void*target,
                                int bytes)
```

`CMFS_physical_transfer_length` takes the arguments of a `CMFS_read_file_physical` or `CMFS_write_file_physical` function call and returns the number of bytes that such a call will transfer. In other words, it returns the same value that a successfully completed call to the read or write function would return.

Because of the complications of padding, we recommend using a physical read only when the file pointer is in the same position it was in when the physical write was performed.

2 Additional Discussion of Return Values of I/O Read and Write Routines

Section 2.4.3 notes that `CMFS_read_file[_always]` and `CMFS_write_file[_always]` return the number of bytes read or written in each position of the parallel variable. In addition, note that, if the requested amount of data isn't actually read, `CMFS_read_file[_always]` returns 0, because it can't sensibly represent the number of bytes read in each position (which may be fractional).